

# Handouts

November 2, 2023 9:20 AM

```
> with(Groebner);
[Basis, FGLM, HilbertDimension, HilbertPolynomial, HilbertSeries, Homogenize, InitialForm,
InterReduce, IsBasis, IsProper, IsZeroDimensional, LeadingCoefficient, LeadingMonomial,
LeadingTerm, MatrixOrder, MaximalIndependentSet, MonomialOrder, MultiplicationMatrix,
MultivariateCyclicVector, NormalForm, NormalSet, RationalUnivariateRepresentation, Reduce,
RememberBasis, SPolynomial, Solve, SuggestVariableOrder, Support, TestOrder, ToricIdealBasis,
TrailingTerm, UnivariatePolynomial, Walk, WeightedDegree]

> f1,f2 := x*y+1, y+1;
f1,f2 := xy + 1, y + 1

> f := x*y^2-x;
f := xy^2 - x

> f1,f2 := x*y+1, y+1;
f1,f2 := xy + 1, y + 1

Divide f by [f1,f2] using lexicographical ordering with x>y
> NormalForm( f, [f1,f2], plex(x,y) );
1 - x
> NormalForm( f, [f2,f1], plex(x,y) );
0
> NormalForm( f, [f1,f2], plex(x,y), 'a' );
1 - x
> a;
[y, -1]
> G := Basis( [f1,f2], plex(x,y) );
G := [y + 1, x - 1]
> NormalForm( f, G, plex(x,y) );
0
> NormalForm( f, [G[2],G[1]], plex(x,y) );
0
```

$$\begin{matrix} f_1 & f_2 \\ \langle xy+1, y+1 \rangle & = \langle -x+1, y+1 \rangle \\ & \downarrow \text{VUL} \\ f_3 & = f_1 - x f_2 = -x+1 \end{matrix}$$

$$\begin{matrix} I = \langle f_1, f_2 \rangle \\ \downarrow \\ I = \langle y+1, x-1 \rangle \end{matrix}$$

## The Division Algorithm for $k[x_1, x_2, \dots, x_n]$ .

**Input**  $\prec$  a monomial ordering on  $\mathbb{Z}_{>0}^n$   
 $f \in k[X]$ , divisors  $f_1, f_2, \dots, f_s \in k[X]$   
 where  $X = x_1, x_2, \dots, x_n, f_i \neq 0$

**Outputs** quotients  $a_1, a_2, \dots, a_s \in k[X]$  and  
 remainder  $r \in k[X]$  satisfying

- (i)  $f = a_1 f_1 + a_2 f_2 + \dots + a_s f_s + r$ .
- (ii)  $LT(f_i)$  does not divide any term in  $r$  and
- (iii)  $LM(f) \geq LM(a_i f_i)$  for  $1 \leq i \leq s$ .

$$(a_1, a_2, \dots, a_s) \leftarrow (0, 0, \dots, 0)$$

$$(r, p) \leftarrow (0, f)$$

**while**  $p \neq 0$  **do**

find the first  $i$  such that  $LT(f_i) | LT(p)$ .

**if**  $\exists i$  **then**  $(r, p) \leftarrow (r + LT(p), p - LT(p))$  **CASE 1**

**else**  $t \leftarrow LT(p)/LT(f_i)$

$(a_i, p) \leftarrow (a_i + t, p - t f_i)$  **CASE 2**

**end if**

**end while**

**output**  $(a_1, a_2, \dots, a_s, r)$

Proof of termination.

I claim each time round the loop either  $p_{\text{new}} = 0$   
 or  $LM(p_{\text{new}}) < LM(p_{\text{old}})$ .

CASE 1.  $LT(f_i) \nmid LT(p)$   $p = LT(p) + \overbrace{(p - LT(p))}^{\neq}$   
 IS  $LM(p - LT(p)) < LM(p)$  or  $p - LT(p) = \underline{0}$ .

CASE 2. IS  $p - t f_i = 0$  or  $LM(p - t f_i) < LM(p)$ ?

Letting  $p_1, p_2, p_3, \dots$  denote the values of  $p$  at the  $i$ th step in the division algorithm we have

$$LM(p_1) > LM(p_2) > LM(p_3) > \dots$$

Since  $>$  is a well ordering Lemma 2 says such a sequence cannot continue indefinitely, hence eventually  $p=0$  and the  $\div$  alg. stops.