

The modular gcd algorithm in  $\mathbb{Z}[x]$  uses  $\phi_p: \mathbb{Z}[x] \rightarrow \mathbb{Z}_p[x]$  and the CRT. What about gcds in  $\mathbb{Z}[x, y, z, \dots]$  and  $\mathbb{F}_q[x, y, z, \dots]$ ??

Eg.  $\gcd(x^4 - t^4, x^6 - t^6)$  in  $\mathbb{Z}[t][x] = 1 \cdot x^2 - t^2$

$(x^2 - t^2)(x^2 + t^2)$  ;  $(x^2 - t^2)(x^4 + t^2x^2 + t^4)$

*unlucky eval. point.*  
 $\phi_{t=0} \gcd(x^4, x^6) = x^4$

$\phi_{t=1} \gcd(x^4 - 1, x^6 - 1) = 1 \cdot x^2 - 1$

$\phi_{t=2} \gcd(x^4 - 2^4, x^6 - 2^6) = 1 \cdot x^2 - 4$

$\phi_{t=3} \gcd(x^4 - 81, x^6 - 3^6) = 1 \cdot x^2 - 9$

$\phi_{t=4} \gcd(x^4 - 256, x^6 - 4^6) = 1 \cdot x^2 - 16$

*Interpolate.*

$\gcd(x^4 + t^4, x^6 + t^6)$  in  $\mathbb{Z}_2[t][x] = x^2 + t^2$

$\phi_{t=0} \gcd(x^4, x^6) = x^4$   $t=0$  is unlucky.

$\phi_{t=1} \gcd(x^4 + 1, x^6 + 1) = x^2 + 1$  ✓

We have only 2 evaluation points in  $\mathbb{Z}_2$ .  
 Where can we get more evaluation points from?